The Complexity of Distance-r Dominating Set Reconfiguration

Niranka Banerjee, RIMS, Kyoto University

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Joint work with Duc A. Hoang



Model: Graph Reconfiguration

Graph Reconfiguration

In a *reconfiguration variant* of a computational problem, two *feasible solutions* S and T are given along with a *reconfiguration rule* that describes how to slightly modify one feasible solution to obtain a new one.

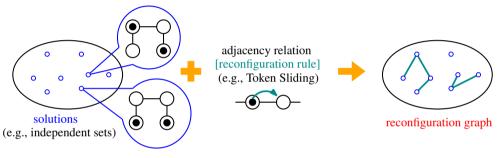


Figure: Reconfiguration

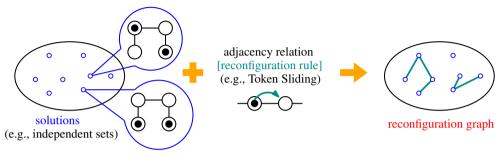


Figure: Reconfiguration

Main questions:

• REACHABILITY: Is there a path between two given solutions? Can we transform S into T via a sequence of feasible solutions.

Such a sequence, if exists, is called a *reconfiguration sequence*.

• SHORTEST PATH: If REACHABILITY is yes, can we find a shortest path between S and T?

Computation problem we consider: Distance-r Dominating Set

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Definition:

Distance-r dominating set (DrDS) of G is a vertex subset D where each vertex of G is within distance r from some member of D.

Distance-r Dominating Set Reconfiguration DrDSR

DrDSR: We solve Reachability of DrDS in the reconfiguration graph.

We prove hardness results of DrDSR in different graph classes for $r \ge 2$ under reconfiguration rule of Token Sliding (TS) and Token Jumping (TJ).

We prove hardness results of DrDSR, $r \ge 2$ in different graph classes.

For example, DrDSR is PSPACE-Complete in Bipartite Graphs, Planar graphs and Chordal Graphs under TS and TJ.

There are other graph classes where DrDSR, $r \ge 2$ is solvable in polynomial time.

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DrDSR can be solved in polynomial time on interval graphs under TJ and co-graphs under both TS and TJ.

Open Questions:

- What is the complexity of DrDSR, $r \ge 2$ under TS on trees?
- What is the complexity of DrDSR, $r \ge 2$ under TS on interval graphs?

Thank You!